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# Robust Algebraic Multilevel Domain Decomposition Preconditioner for General Sparse Matrices 

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STFC, Rutherford Appleton Laboratory
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## Collaborators

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## Outline

Introduction to (Algebraic) Additive Schwarz

Adaptive Coarse Spaces
Some Existing CS
New CS
Normal Eqs.
More CS

Summary

## Motivation

Solving sparse linear systems is omnipresent in scientific computing.

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A x=b
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$A \in \mathbb{K}^{n \times n}$ very large, sparse, and ill conditioned.

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- Black box: prefer non-intrusive solvers. Only provide $A$ and $b$
- Easy setup: Few knowledge on linear solvers. Not worry how to set it up perfectly. As minimal parameters as possible.
(Type: Hermitian, saddle-point, etc; Accuracy; Max iter)


## Overlapping DD

Solve the Poisson equation in $\Omega$


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$$
R_{1} A R_{1}^{T} x=R_{1} b
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## Overlapping DD

Solve the Poisson equation in $\Omega$


$$
R_{2} A R_{2}^{\top} x=R_{2} b
$$

## Overlapping DD

Solve the Poisson equation in $\Omega$


$$
R_{1}^{T}\left(R_{1} A R_{1}^{T}\right)^{-1} R_{1}+R_{2}^{T}\left(R_{2} A R_{2}^{T}\right)^{-1} R_{2}
$$

Iterate updating the solution values.

## Overlapping DD



Figure: Semistructured mesh decomposed into 32 overlapping subdomains
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## Ingredients: Overlapping Subdomain

The sparsity graph of $A$ has $n$ nodes.
$N$ non-overlapping sudbomains $\left\{\Omega_{l i}\right\}_{1 \leq i \leq N}: N$ disjoint subsets of $\Omega=\llbracket 1, n \rrbracket . n_{l i}=\# \Omega_{l i}$

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\end{array}\right), \quad N=2
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\sum_{i=1}^{N} R_{i}^{T} D_{i} R_{i}=I_{n}
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- $R_{/ 1}=I\left(\left[\begin{array}{ll}1 & 2\end{array}\right],:\right), R_{\Gamma 1}=I(3,:), R_{1}=I\left(\left[\begin{array}{lll}1 & 2 & 3\end{array}\right],:\right)$
$-R_{I 2}=I\left(\left[\begin{array}{ll}3 & 4\end{array}\right],:\right), R_{\Gamma 2}=I(2,:), R_{2}=I\left(\left[\begin{array}{lll}3 & 4 & 2\end{array}\right],:\right)$


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$$
D_{1}=\left(\begin{array}{ccc}
1 & & \\
& 1 & \\
& & 0
\end{array}\right) \quad D_{2}=\left(\begin{array}{lll}
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a_{33} & a_{34} & a_{32} \\
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a_{23} & & a_{22}
\end{array}\right)
$$

Define $P_{i}=I\left(\left[R_{l i}, R_{\Gamma i}, R_{c i}\right],:\right)$,

$$
P_{i} A P_{i}^{\top}=\left(\begin{array}{lll}
A_{l i, l i} & A_{l i, \Gamma i} & \\
A_{\Gamma i, l i} & A_{\Gamma i, \Gamma i} & A_{\Gamma i, c i} \\
& A_{c i, \Gamma i} & A_{c i, c i}
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$$

## $\delta$-Overlap

Through the sparsity graph, define $\Omega_{\Gamma_{1: \delta} i}$ Define

$$
P_{i}=I\left(\left[R_{l i}, R_{\Gamma_{1: \delta-1} i}, R_{\Gamma_{\delta} i}, R_{c i}\right],:\right),
$$

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$R_{i}=\left[R_{l i}, R_{\Gamma_{1: \delta-1} i}, R_{\Gamma_{\delta}}\right]$.

## One-Level Schwarz

Four stages:

1. Restrict
2. Solve locally
3. Augment
4. Update

$$
M_{1}^{-1}=\sum_{i=1}^{N} R_{i}^{T} A_{i i}^{-1} R_{i}
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## One-Level Schwarz Not Scalable

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M_{1}^{-1}=\sum_{i=1}^{N} R_{i}^{T} A_{i i}^{-1} R_{i}
$$

| $N$ | 2 | 4 | 8 | 16 | 32 | 64 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| It | 42 | 53 | 66 | 74 | 84 | 97 |

Table: 2D Poisson on $300 \times 300$ mesh. Metis partitioning.

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Iteration count $=f(N)$. Not scalable

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Table: 2D Poisson on $300 \times 300$ mesh. Metis partitioning.

Iteration count $=f(N)$. Not scalable
Need a second level (coarse space correction) to maintain robustness

$$
M_{2}^{-1}=R_{0}^{H} A_{00}^{-1} R_{0}+\sum_{i=1}^{N} R_{i}^{T} A_{i i}^{-1} R_{i}
$$

## Adaptive Coarse Spaces (for Overlapping Schwarz) I

PDE based (Two-level)

- A coarse space construction based on local Dirichlet-to-Neumann maps [Nataf et al., 2011]
- Abstract robust coarse spaces for systems of PDEs via generalized eigenproblems in the overlaps [Spillane et al., 2014]
- SHEM: an optimal coarse space for RAS and Its multiscale approximation [Gander and Loneland, 2017]
- Adaptive GDSW coarse spaces of reduced dimension for overlapping Schwarz methods [Heinlein et al., 2020]
- A multilevel Schwarz preconditioner based on a hierarchy of robust coarse spaces [Al Daas et al., 2021]
- A comparison of coarse spaces for Helmholtz problems in the high frequency regime [Bootland et al., 2021]
- Multilevel spectral domain decomposition [Bastian et al., 2022]
- A fully algebraic and robust two-level Schwarz method based on optimal local approximation spaces [Heinlein and Smetana, 2022]


## Adaptive Coarse Spaces (for Overlapping Schwarz) II

Fully Algebraic

- A class of efficient locally constructed preconditioners based on coarse spaces [Al Daas and Grigori, 2019]
- Fully algebraic domain decomposition preconditioners with adaptive spectral bounds [Gouarin and Spillane, 2021]
- A Robust Algebraic Domain Decomposition Preconditioner for Sparse Normal Equations [Al Daas et al., 2022b]
- A robust algebraic multilevel domain decomposition preconditioner for sparse symmetric positive definite matrices [AI Daas and Jolivet, 2022]
- Efficient algebraic two-level Schwarz preconditioner for sparse matrices [Al Daas et al., 2022a]


## Interface-to-Interior Operator

$$
\begin{aligned}
& P_{i}=I\left(\left[R_{l i}, R_{\Gamma_{1: \delta-1} i}, R_{\Gamma_{\delta i} i}, R_{c i}\right],:\right), \\
& \quad P_{i} A P_{i}^{\top}=\left(\begin{array}{cccc}
A_{l i, l i} & A_{l i, \Gamma_{1: \delta-1} i} & \\
A_{\Gamma_{1: \delta-1} i, l i} & A_{\Gamma_{1: \delta-1} i, \Gamma_{1: \delta-1} i} & A_{\Gamma_{1: \delta-1} i, \Gamma_{\delta} i} & \\
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\end{array}\right)
\end{aligned}
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& & A_{c i, \Gamma_{\delta} i} & A_{c i, c i}
\end{array}\right)
$$

Consider $T_{i}: x_{\delta i} \in \Omega_{\Gamma_{\delta} i} \mapsto x_{l i} \in \Omega_{l i}$, the restriction of the corresponding local solution.

$$
\left(\begin{array}{ccc}
A_{l i, l i} & A_{l i, \Gamma_{1: \delta-1} i} & \\
A_{\Gamma_{1: \delta-1} i, l i} & A_{\Gamma_{1: \delta-1} i, \Gamma_{1: \delta-1} i} & A_{\Gamma_{1: \delta-1} i, \Gamma_{\delta i} i}
\end{array}\right)\left(\begin{array}{c}
x_{l i} \\
x_{\Gamma_{1: \delta-1} i} \\
x_{\delta i}
\end{array}\right)=\left(\begin{array}{c}
0 \\
0 \\
x_{\delta i}
\end{array}\right)
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A_{\Gamma_{1: \delta-1} i, l i} & A_{\Gamma_{1: \delta-1} i, \Gamma_{1: \delta-1} i} & A_{\Gamma_{\delta} i, c i} \\
& A_{\Gamma_{\delta} i, \Gamma_{1: \delta-1} i} & A_{\Gamma_{\delta} i, \Gamma_{\delta} i} \\
& A_{c i, \Gamma_{\delta} i} & A_{c i, c i}
\end{array}\right) \\
& \left.T_{i}\left(x_{\delta i}\right)=x_{l i}=\left(\begin{array}{cc}
A_{l i, l i} & A_{l i, \Gamma_{1: \delta-1} i} \\
A_{\Gamma_{1: \delta-1} i, l i} & A_{\Gamma_{1: \delta-1} i, \Gamma_{1: \delta-1} i}
\end{array}\right)\binom{0}{-A_{\Gamma_{1: \delta-1} i, \Gamma_{\delta} i} x_{\delta i}}\right)_{l i}
\end{aligned}
$$

## Theorem

To be submitted [HAD, Jolivet, Nataf, Tournier]

## Theorem

Set $R_{0}^{H}=\left[R_{I 1} Z_{1}, \ldots, R_{I N} Z_{N}\right]$, where $Z_{i}=\operatorname{Im}\left(T_{i}\right)$

$$
M_{2}^{-1}=R_{0}^{H} A_{00}^{-1} R_{0}+\sum_{i=1}^{N} R_{i}^{T} A_{i i}^{-1} R_{i}
$$

If $A$ is HPD:

$$
\kappa\left(M_{2}^{-1} A\right)=C
$$

C depends only on the largest number of neighbouring overlapping subdomains

## SVD Interface-to-Interior Operator



Figure: $N=32$. Subdomain 1. Singular values of the interface-to-interior operator for a Poisson equation.

## SVD Interface-to-Interior Operator



Figure: $N=32$. Subdomain 1. Singular values of the interface-to-interior operator for the matrix $A^{T} A$, where $A$ is the Ruccil matrix (SSMC).

## Theorem

To be submitted [HAD, Jolivet, Nataf, Tournier]

## Theorem

Set $R_{0}^{H}=\left[R_{/ 1} Z_{1}, \ldots, R_{I N} Z_{N}\right]$, where $Z_{i}=t S V D\left(\operatorname{Im}\left(T_{i}\right), \varepsilon\right)$

$$
M_{2}^{-1}=R_{0}^{H} A_{00}^{-1} R_{0}+\sum_{i=1}^{N} R_{i}^{T} A_{i i}^{-1} R_{i}
$$

If $A$ is HPD:

$$
\kappa\left(M_{2}^{-1} A\right)=C\left(1+\kappa\left(M_{1}^{-1} A\right) \varepsilon\right)
$$

C depends only on the largest number of neighbouring overlapping subdomains

## Numerical experiments

| $N$ | 4 | 8 | 16 | 32 |
| :--- | :---: | :---: | :---: | :---: |
| Diffusion | $14(160)$ | $12(320)$ | $11(640)$ | $8(1280)$ |
| Adv-Diff | $14(160)$ | $13(320)$ | $14(640)$ | $12(1280)$ |
| Stokes | $47(320)$ | $42(640)$ | $43(1280)$ | $49(2559)$ |
| Biharmonic | $51(240)$ | $55(480)$ | $34(960)$ | $22(1920)$ |
| Elasticity | $50(320)$ | $37(640)$ | $36(1267)$ | $28(2529)$ |

Table: Strong scaling on variety of problems

## Numerical experiments

| $N$ | 8 | 32 |
| :--- | :---: | :---: |
| Diffusion | $10(160)$ | $13(640)$ |
| Adv-Diff | $12(160)$ | $13(640)$ |
| Stokes | $32(614)$ | $49(2555)$ |
| Biharmonic | $11(639)$ | $15(2560)$ |
| Elasticity | $18(554)$ | $28(2529)$ |

Table: Weak scaling on variety of problems

## Nornal Equations

[HAD, Jolivet, Scott. SISC 22']

## Theorem

For a sparse $A$ with $A=B^{H} B$, or $A=B^{H} \operatorname{diag}(g) B, g \geq 0$

$$
\kappa\left(M_{2}^{-1} A\right)=C(1+\tau)
$$

where $\tau>0$ is a user-specified.

## PDE-CO

Solve

$$
\min _{y}\|y-\hat{y}\|_{\Omega_{1}}^{2}+\beta\|u\|_{\Omega_{2}}^{2} \quad \text { subject to } \mathcal{L} y=u \text { in } \Omega
$$

The resulting matrix

$$
\left(\begin{array}{ccc}
M & & K^{*} \\
& \beta R & L^{*} \\
K & L &
\end{array}\right)
$$

Mass lumping yields an equivalent diagonal matrix $W$ to the (1:2,1:2)-block. $\widetilde{S}=J^{*} J$, where $J^{*}=[K L] W^{-1 / 2}$.

## Poisson PDE-CO

IFISS: Grid $2^{8} \times 2^{8}, \beta=0.01, Q_{2}-F E$, matrix length $\approx 200 K$.


Figure: Residual history

## Helmholtz PDE-CO

Test case inspired by Kouri et al. 21'
Grid $160 \times 160, \beta=10^{-5}, P_{1}$-FE, matrix length $\approx 50 K$.



Figure: State (real part): Desired (left), solution (right)

## Helmholtz PDE-CO

Test case inspired by Kouri et al. 21'
Grid $160 \times 160, \beta=10^{-5}, P_{1}$-FE, matrix length $\approx 50 \mathrm{~K}$.


Figure: 3D view of the state (real part)

## Helmholtz PDE-CO

Test case inspired by Kouri et al. 21'
Grid $160 \times 160, \beta=10^{-5}, P_{1}-$ FE, matrix length $\approx 50 K$.


Figure: Control (real part)

## Helmholtz PDE-CO

Test case inspired by Kouri et al. 21'
Grid $160 \times 160, \beta=10^{-5}, P_{1}-$ FE, matrix length $\approx 50 K$.


Figure: Residual history

## Diagonally Dominant HPD

[HAD, Jolivet, Rees. SISC 23']

## Theorem

For a sparse A HPD

$$
\kappa\left(M_{2}^{-1} A\right)=C(1+\tau)
$$

where $\tau>0$ is a user-specified.

## Highly Non Symmetric

$$
\nabla \cdot(V u)-\nu \nabla \cdot(\kappa \nabla u)=0 \text { in } \Omega \quad u=0 \text { on } \Gamma_{0} \quad u=1 \text { on } \Gamma_{1}
$$


(a) Mesh

(b) $\nu=10^{-2}$

(c) $\nu=10^{-4}$

## Highly Non Symmetric

$$
\nabla \cdot(V u)-\nu \nabla \cdot(\kappa \nabla u)=0 \text { in } \Omega \quad u=0 \text { on } \Gamma_{0} \quad u=1 \text { on } \Gamma_{1}
$$


(a) Mesh

(b) $\nu=10^{-2}$

(c) $\nu=10^{-4}$


$$
0
$$

| Prec | Dimension | $k$ | $N$ | $n$ | 1 | $10^{-1}$ | $10^{-2}$ | $10^{-3}$ | $10^{-4}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $M_{2}^{-1}$ | 2 | 1 | 1,024 | $6.3 \cdot 10^{6}$ | 23 (52,875) | $20(52,872)$ | $19_{(52,759)}$ | $20(47,497)$ | $21(28,235)$ |
|  | 3 | 2 | 4,096 | $8.1 \cdot 10^{6}$ | 18 (1.8 • $10^{5}$ ) | $14\left(1.8 \cdot 10^{5}\right)$ | 11 (1.6-10 ${ }^{5}$ ) | $16{ }_{(97,657)}$ | $29(76,853)$ |
| GAMG | 23and3 <br> $y$ <br> Council3 | 1 | 1,024 | $6.3 \cdot 10^{6}$ | 42 | 48 | 88 | $\dagger$ | $\dagger$ |
|  |  | 2 | 4,096 | $8.1 \cdot 10^{6}$ | 40 | 38 | 65 | $\dagger$ | $\dagger$ |
|  |  | 1 | 1,024 | $6.3 \cdot 10^{6}$ | 50 | 49 | 19 | 7 | $\dagger$ |
|  |  | 2 | 4,096 | $8.1 \cdot 10^{6}$ | 12 | 9 | 7 | $\dagger$ | $\dagger$ |

## Summary \& Perspectives

Summary:

- Algebraic DD provides a simple way to construct preconditioners that are effective, efficient, black-box and easy to set up
- Provable: Diagonally weighted normal equations matrix (Schur complement); HPD; Diagonally dominant HPD
- All preconditioner are accessible in PETSc PCHPDDM

Perspectives:

- Extension to general Schur complement $B^{H} G^{-1} B$
- Extend theory to non-Hermitian matrices


## Thank you for your attention!

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$$
\begin{gathered}
\nabla \cdot(V u)-\nu \nabla \cdot(\kappa \nabla u)=0 \text { in } \Omega \quad u=0 \text { on } \Gamma_{0} \quad u=1 \text { on } \Gamma_{1} \\
v(x, y)=\binom{x(1-x)(2 y-1)}{-y(1-y)(2 x-1)} \quad \text { or } \quad v(x, y, z)=\left(\begin{array}{c}
2 x(1-x)(2 y-1) z \\
-y(1-y)(2 x-1) \\
-(1-z)(2 x-1)(2 y-1)
\end{array}\right),
\end{gathered}
$$

## Two to multi-level [H.A., P.J., L.G.,P.-H.T. SISC '21]

$$
A_{00}=R_{0} A R_{0}^{H}
$$

## Two to multi-level [H.A., P.J., L.G.,P.-H.T. SISC '21]

$$
\begin{gathered}
A_{00}=R_{0} A R_{0}^{H} \\
u^{H} \sum_{i=1}^{N} \widetilde{A}_{i} u \leq k_{m} u^{H} A u
\end{gathered}
$$

## Two to multi-level [H.A., P.J., L.G.,P.-H.T. SISC '21]

$$
\begin{gathered}
A_{00}=R_{0} A R_{0}^{H} \\
u^{H} \sum_{i=1}^{N} \widetilde{A}_{i} u \leq k_{m} u^{H} A u \\
\left(R_{0} v\right)^{H} \sum_{i=1}^{N} \widetilde{A}_{i}\left(R_{0} v\right) \leq k_{m}\left(R_{0} v\right)^{H} A\left(R_{0} v\right)
\end{gathered}
$$

## Two to multi-level [H.A., P.J., L.G.,P.-H.T. SISC '21]

$$
\begin{gathered}
A_{00}=R_{0} A R_{0}^{H} \\
u^{H} \sum_{i=1}^{N} \tilde{A}_{i} u \leq k_{m} u^{H} A u \\
\left(R_{0} v\right)^{H} \sum_{i=1}^{N} \widetilde{A}_{i}\left(R_{0} v\right) \leq k_{m}\left(R_{0} v\right)^{H} A\left(R_{0} v\right) \\
v^{H} \sum_{i=1}^{N}\left(R_{0}^{H} \widetilde{A}_{i} R_{0}\right) v \leq k_{m} v^{H}\left(R_{0}^{H} A R_{0}\right) v=k_{m} v^{H} A_{00} v
\end{gathered}
$$

## Two to multi-level [H.A., P.J., L.G.,P.-H.T. SISC '21]

$$
\begin{gathered}
A_{00}=R_{0} A R_{0}^{H} \\
u^{H} \sum_{i=1}^{N} \widetilde{A}_{i} u \leq k_{m} u^{H} A u \\
\left(R_{0} v\right)^{H} \sum_{i=1}^{N} \widetilde{A}_{i}\left(R_{0} v\right) \leq k_{m}\left(R_{0} v\right)^{H} A\left(R_{0} v\right) \\
v^{H} \sum_{i=1}^{N}\left(R_{0}^{H} \widetilde{A}_{i} R_{0}\right) v \leq k_{m} v^{H}\left(R_{0}^{H} A R_{0}\right) v=k_{m} v^{H} A_{00} v \\
v^{H} \sum_{\substack{H=1}}^{N_{2}} \underbrace{\left(\sum_{i \in G_{j}}\left(R_{0}^{H} \widetilde{A}_{i} R_{0}\right)\right)}_{\substack{\text { Science and } \\
\text { Technoligy } \\
\text { Facilities council }}} v \leq k_{m} v^{H}\left(R_{0}^{H} A R_{0}\right) v=k_{m} v^{H} A_{00} v \\
\widetilde{A}_{00, j}
\end{gathered}
$$

## GenEO FEM [Dolean et al '15]

$$
\begin{aligned}
a(u, v) & =\sum_{K \in \mathcal{T}} \int_{K} u v \rightarrow A \\
\widetilde{a}(u, v) & =\sum_{K \in \mathcal{T}_{i}} \int_{K} u v \rightarrow \widetilde{A}_{i}
\end{aligned}
$$

