## An Explanation of the MRRR algorithm to compute eigenvectors of symmetric tridiagonal matrices

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## Difficulties

- All eigenvalues of $T$ are easily computed in $O\left(n^{2}\right)$ time.
- Given $\hat{\lambda}$, inverse iteration computes the eigenvector:

$$
(T-\hat{\lambda} I) x_{i+1}=x_{i}, \quad i=0,1,2, \ldots
$$

- Costs $O(n)$ per iteration.
- Typically, 1-3 iterations are enough.
- BUT, inverse iteration only guarantees

$$
\|T \hat{v}-\hat{\lambda} \hat{v}\|=O(\varepsilon\|T\|)
$$

## Fundamental Limitations

Gap Theorem :

$$
\sin \angle(v, \hat{v}) \leq \frac{\|T \hat{v}-\hat{\lambda} \hat{v}\|}{\operatorname{Gap}(\hat{\lambda})}
$$

Assume all off-diagonals not negligible, so $\operatorname{Gap}(\hat{\lambda})$ not zero, but can be small:

$$
\left[\begin{array}{cccc}
1 & \varepsilon_{1} & & \\
\varepsilon_{1} & 1 & \varepsilon_{2} & \\
& \varepsilon_{2} & 1 & \varepsilon_{3} \\
& & \varepsilon_{3} & 1
\end{array}\right]
$$

When eigenvalues are close, independently computed eigenvectors WILL NOT be mutually orthogonal.

## Eigenvalues of Biphenyl Matrix

- Plot eigenvalues :

- Plot $\operatorname{Absgap}(i)=\log _{10}\left(\min \left(\lambda_{i+1}-\lambda_{i}, \lambda_{i}-\lambda_{i-1}\right) /\|T\|\right)$ :

- LAPACK — one big "cluster" $\lambda_{1}, \lambda_{2}, \ldots, \lambda_{939}$.
- Tridiagonal solution takes $80 \%$ of Total Time.


## Ideal Solution

- Fundamental Limitation:

$$
\sin \angle(v, \hat{v}) \leq \frac{\|T \hat{v}-\hat{\lambda} \hat{v}\|}{\operatorname{Gap}(\hat{\lambda})}
$$

- Get smallest possible residual norm.
- Compute eigenvalue to greatest accuracy possible :

$$
|\hat{\lambda}-\lambda|=O(\varepsilon|\hat{\lambda}|)
$$

- Compute eigenvector to high relative accuracy :

$$
\|T \hat{v}-\hat{\lambda} \hat{v}\|=O(\varepsilon|\hat{\lambda}|)
$$

- Gap Theorem implies :

$$
\sin \angle(v, \hat{v})=\frac{O(\varepsilon|\lambda|)}{\operatorname{Gap}(\hat{\lambda})}=\frac{O(\varepsilon)}{\operatorname{Relgap}(\hat{\lambda})}
$$

- Can we achieve the above ?

Key Idea 1. Discard Tridiagonals, Embrace Bidiagonals

## Factored Forms yield Better Representations

- Tridiagonals DO NOT determine their eigenvalues to high relative accuracy.
- Bidiagonals determine their singular values to high relative accuracy.

$$
\begin{aligned}
& T+\mu I \quad L \quad L^{T}
\end{aligned}
$$

- Bidiagonal Factors are "better" since they allow us to
- compute eigenvalues to high relative accuracy,
- compute eigenvectors to high relative accuracy.
- High accuracy $\Rightarrow$ Orthogonality.
- For interior eigenvalues, extends to indefinite factorization $L D L^{T}$.


## Algorithm Outline

1. Choose $\mu$ such that $T+\mu I$ is positive definite.
2. Compute the factorization :

$$
T+\mu I=L D L^{T}
$$

3. Compute eigenvalues of $L D L^{T}$ to high relative accuracy (by dqds or bisection).
4. Given eigenvalues, compute accurate eigenvectors of $L D L^{T}$.

## - HOW?

## Key Idea 2. Shift with Differential QD

How do we get an eigenvector such that

$$
\|T \hat{v}-\hat{\lambda} \hat{v}\|=O(\varepsilon|\hat{\lambda}|) ?
$$

## Differential Transformations

- Inverse iteration - Solve for $z$ :

$$
\begin{aligned}
L D L^{T}-\hat{\lambda} I & =L_{+} D_{+} L_{+}^{T} \\
L_{+} D_{+} L_{+}^{T} z & =\text { random vector. }
\end{aligned}
$$

Simple qd : $\quad D_{+}(1):=d_{1}-\hat{\lambda}$

$$
\text { for } i=1, n-1
$$

$$
L_{+}(i):=\left(d_{i} l_{i}\right) / D_{+}(i)
$$

$$
D_{+}(i+1):=d_{i} l_{i}^{2}+d_{i+1}-L_{+}(i) d_{i} l_{i}-\hat{\lambda}
$$

end for

| Differential qd (dqds) $:$ | $s_{1}:=-\hat{\lambda}$ |
| :---: | :--- |
|  | for $i=1, n-1$ |
| $D_{+}(i):=s_{i}+d_{i}$ |  |
| $L_{+}(i):=\left(d_{i} l_{i}\right) / D_{+}(i)$ |  |
| $s_{i+1}:=L_{+}(i) l_{i} s_{i}-\hat{\lambda}$ |  |
| end for |  |
|  | $D_{+}(n):=s_{n}+d_{n}$ |

## Computing an Eigenvector

- Compute the appropriate Twisted Factorization :

$$
T-\hat{\lambda} I=N_{r} D_{r} N_{r}^{T},
$$


and $r$ is chosen to minimize $\left|\gamma_{r}\right|$ (it will be $O(\lambda-\hat{\lambda})$ ).

- Solve for $z, N_{r} D_{r} N_{r}^{T} z=\gamma_{r} e_{r}\left(\Rightarrow N_{r}^{T} z=e_{r}\right)$ :

$$
z(i)= \begin{cases}1, & i=r \\ -L_{+}(i) \cdot z(i+1), & i=r-1, \ldots, 1, \\ -U_{-}(i-1) \cdot z(i-1), & i=r+1, \ldots, n\end{cases}
$$

- Solves an open problem posed by Wilkinson (1965).


## Main Theorem

THEOREM. [Dhillon \& Parlett, 2003] Eigenvectors computed by twisted factorization are numerically orthogonal if eigenvalues of $L D L^{T}$ have large relative gaps. In particular,

$$
\left(\hat{v}_{i}, \hat{v}_{j}\right)=\frac{O(\varepsilon)}{\operatorname{Relsep}\left(\lambda_{i}, \lambda_{j}\right)},
$$

where

$$
\operatorname{Relsep}\left(\lambda_{i}, \lambda_{j}\right)=\frac{\left|\lambda_{i}-\lambda_{j}\right|}{\max \left(\left|\lambda_{i}\right|,\left|\lambda_{j}\right|\right)}
$$

- Example of Large Relsep :

$$
\lambda_{1}=10^{-16}, \lambda_{2}=10^{-15} \Rightarrow \operatorname{Relsep}\left(\lambda_{1}, \lambda_{2}\right) \approx 1
$$

Above Theorem $\Rightarrow$ Automatic Orthogonality.

- Example of Small Relsep :

$$
\begin{aligned}
& \lambda_{1}=1.000000000000001, \\
& \lambda_{2}=1.000000000000002 .
\end{aligned}
$$

## Proof of Correctness

- Desired Relationship: $L D L^{T}-\hat{\lambda} I=N_{r} D_{r} N_{r}^{T}$, and $N_{r} D_{r} N_{r}^{T} z=\gamma_{r} e_{r}$.

- Exact Mathematical relationship holds : $\bar{L} \bar{D} \bar{L}^{T}-\hat{\lambda} I=\tilde{N}_{r} \tilde{D}_{r} \tilde{N}_{r}^{T}$.
- Key step in proof is to relate $\hat{z}$ to $v$ in 3 steps :

1. $\hat{z}$ is close to $\tilde{z}$,
(only multiplications),
2. $\sin \angle(\bar{v}, \tilde{z})=O(\varepsilon|\bar{\lambda}|) / \operatorname{gap}(\hat{\lambda}), \quad\left(\left|\tilde{\gamma}_{r}\right|=O(\varepsilon|\bar{\lambda}|)\right)$,
3. $\sin \angle(\bar{v}, v)=O(\varepsilon) / \operatorname{relgap}(\hat{\lambda}) \quad$ (relative perturbation theory).

$$
\Rightarrow \quad \sin \angle(\hat{z}, v)=\frac{O(\varepsilon)}{\operatorname{Relgap}(\hat{\lambda})} .
$$

Key Idea 3. Shift for Separation, again differentially

## Algorithm $\mathrm{MR}^{3}$ (Multiple RRRs)

1. Choose $\mu$ such that $T+\mu I$ is positive definite.
2. Compute the factorization:

$$
T+\mu I=L D L^{T} .
$$

3. Compute eigenvalues of $L D L^{T}$ to high relative accuracy (by dqds or bisection).
4. Group eigenvalues according to their Relative Gaps :
a) isolated (agree in $<3$ digits). Compute eigenvector using a twisted factorization.
b) clustered (agree in $>3$ digits).

- Pick $\mu$ near cluster to form $L D L^{T}-\mu I=L_{1} D_{1} L_{1}^{T}$ (by dqds).
- "Refine" eigenvalues in cluster to high relative accuracy.
- Set $L \leftarrow L_{1}, D \leftarrow D_{1}$. Repeat step 4 for eigenvalues in cluster.


## Key Idea 4. Analyze the Representation Tree

Step 4 of the $\mathbf{M R}^{3}$ algorithm may be represented as a tree:

- At the root is the original factorization.
- At each internal node is a factorization for another shift $\mu$.
- Each child of the node for $\mu$ corresponds to an isolated eigenvalue or a cluster.


## Small Example

- Eigenvalues: $\varepsilon, 1+\sqrt{\varepsilon}, 1+2 \sqrt{\varepsilon}, 2$.
- Extra representation needed at $\sigma=1$ :

$$
L D L^{T}-I=L_{1} D_{1} L_{1}^{T} .
$$

- The following Representation Tree captures the steps of the algorithm:



## Wilkinson's Matrix

- $W_{21}^{+}: 21 \times 21$ Wilkinson's matrix.
- $\lambda_{20}$ and $\lambda_{21}$ are identical to working precision.
- What happens in this case?

$$
L D L^{T}-\hat{\lambda}_{21} I=L_{1} D_{1} L_{1}^{T}
$$

$$
\begin{aligned}
\lambda_{20}\left(L_{1} D_{1} L_{1}^{T}\right) & \& \lambda_{21}\left(L_{1} D_{1} L_{1}^{T}\right) \quad \text { - no digits in common! } \\
-7.28 \times 10^{-14} & \&-1.22 \times 10^{-15} \\
\left(\hat{v}_{20}, \hat{v}_{21}\right) & =1.0 \times 10^{-16}
\end{aligned}
$$

- Computed Eigenvectors $\hat{v}_{20}$ and $\hat{v}_{21}$ :



## Worst Case / Large Depth

$13 \times 13$ matrix with eigenvalues: $0,1,1 \pm 10^{-15}, 1 \pm 10^{-12}, 1 \pm 10^{-9}, 1 \pm 10^{-6}, 1 \pm 10^{-3}, 2$.


## Performance of MRRR on Biphenyl Matrix

For the biphenyl matrix ( $n=966$ ),

- the root node had 805 leaf children and 63 internal node children.
- all nodes at the next level were leaf nodes.
- 49 clusters had 2 eigenvalues,
- 13 clusters had 3-8 eigenvalues.
- one cluster had 9 eigenvalues,


## Residual Norms for computed $z$ - Part 1

- Paper 1 guarantees small residuals at the bottom of the representation tree :
a leaf and its parent

$$
\|(\text { leaf }-\delta \lambda I) z \|=|\gamma|=O(\varepsilon \delta \lambda)
$$

- Our problem :

$$
\|(\text { root }-\lambda I) z \|=|\gamma|=O(\varepsilon \text { spdiam }(\text { root })) \quad ? ? ?
$$

- In exact arithmetic, root residual is also $O(\varepsilon \delta \lambda)$.


## Residual Norms for computed $z$ - Part 2

Compare child residual with parent residual at each interval node from leaf to root.


- By design

$$
\tilde{r}_{p}=\bar{r}_{c} \quad \text { exact }
$$

- also

$$
\begin{gathered}
\tilde{r}_{p}=r_{p}+\delta T_{p} z, \quad \tilde{T}_{p}=T_{p}+\delta T_{p} \\
\bar{r}_{c}=r_{c}+\delta T_{c} z, \quad \bar{T}_{c}=T_{c}+\delta T_{c} \\
T+\delta T=(L+\delta L)(D+\delta D)(L+\delta L)^{T}, \quad T=L D L^{T}
\end{gathered}
$$

## Technical Lemma

If

$$
\begin{aligned}
\|D z\| & \leq c \operatorname{spdiam}\left(T_{0}\right) \\
\left\|\dot{L} D \dot{L}^{T} z\right\| & \leq c \operatorname{spdiam}\left(T_{0}\right), \quad \stackrel{L}{L}=L-I,
\end{aligned}
$$

then

$$
(*) \quad\|\delta T z\| \leq\left(2 c+\frac{1}{2}\right)(9 \varepsilon) \operatorname{spdiam}\left(T_{0}\right)+O\left(n \varepsilon^{2}\right)
$$

NOTE: large values in $D$ and $L D L^{T}$ can be neutralized by small entries in $z$.
$\left.{ }^{*}\right)$ gives a bound on increase in residual norm at each internal node on path from leaf to root.

## Orthogonality



- By Paper 1, eigenvectors with same parent are orthogonal to working accuracy.
- Our problem: $\quad\left|z_{j}^{T} z_{k}\right|=O(n \varepsilon) \quad$ ????
- $\mathcal{S}_{\Gamma}=\mathcal{S}_{\Gamma}^{L D L^{T}}=$ subspace invariant under $L D L^{T}$ for eigenvalues in $\Gamma$.


## Two Angles

$$
\begin{aligned}
\Psi_{k, \Gamma} & :=\angle\left(z_{k}, \mathcal{S}_{\Gamma}\right) \\
\Phi_{\Gamma_{\alpha}} & :=\angle\left(\mathcal{S}_{\Gamma_{\alpha}}^{\text {parent }}, \mathcal{S}_{\Gamma_{\alpha}}^{\text {child }}\right), \quad \mathcal{S}_{\Gamma_{\alpha}}^{\text {parent }} \subset \mathcal{S}_{\Gamma}^{\text {parent }}
\end{aligned}
$$

Lemma 1. $\quad \sin \Psi_{j, \Gamma} \leq \sin \Psi_{j, \Gamma_{\alpha}}+\sin \Phi_{\Gamma_{\alpha}}$
Lemma 2. $\sin \Phi_{\Gamma \alpha} \leq R n \varepsilon$,
$R$ depends on tolerance for relgap.

## Theorem

Let $\left(L D L^{T}, \Gamma\right)$ be the least common ancestor of $z_{j}$ and $z_{k}, j \neq k$. If all internal nodes on the paths from leaves $\langle j\rangle$ and $<k>$ to $\Gamma$, (in the representation tree) are RRRs, then $\cos \angle\left(z_{j}, z_{k}\right) \leq 2$ leafbound $+\{\operatorname{depth}(\Gamma, j)+\operatorname{depth}(\Gamma, k)-2\} R n \varepsilon$.

