# Too large to handle directly? 

Jennifer Scott<br>Joint work with John Reid

## Sparse systems

We are interested in solving

$$
\mathbf{A x}=\mathbf{b}
$$

where $\mathbf{A}$ is

## LARGE

$\begin{array}{llllll}s & p & a & s & e\end{array}$

- Problem sizes of order $>10^{7}$ not uncommon and growing larger
- Direct methods (eg $\left.\mathbf{A}=(\mathbf{P L}) \mathbf{D}(\mathbf{P L})^{\mathbf{T}}\right)$ are popular because they are robust
- But their storage requirements generally grow rapidly with problem size
- One possible solution: use an out-of-core direct solver


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- One possible solution: use an out-of-core direct solver

An out-of-core solver holds the matrix factors in files and may also hold the matrix data and some work arrays in files.

## Brief history of HSL out-of-core solvers

- MA32 frontal solver for element problems, written by Iain Duff in 1980.
- Optionally used direct access files to hold the matrix factors.
- Extended 1983 to assembled unsymmetric systems.
- Superseded in 1992 by MA 42 (Duff and Scott). Major change: use of level 3 BLAS.
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TREESOLV is really the inspiration for our recent codes.

## Other out-of-core solvers

- BCSEXT-LIB (Boeing)
- Oblio (Dobrian and Pothen)
- TAUCS (Toledo and students)
- MUMPS parallel solver: now offers out-of-core version
- Also work by Rothberg and Schreiber

HSL_MA77 is our new out-of-core solver

- HSL_MA77 is designed to solve LARGE sparse symmetric systems, both positive definite and indefinite
- HSL_MA77 implements a multifrontal algorithm
- Matrix data, matrix factor, and the main work space (multifrontal stack) held in files

Aim today: to provide brief introduction to HSL_MA77 and to present some numerical results .... hope you will go away wanting to try the code

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- Reverse communication interface with input by rows or by elements
- Separate calls for each phase
- Entering of integer and real matrix data
- Analyse phase (set up data structures using user-supplied pivot order)
- Factorization (compute and store factor plus optional solve)
- Solve (any number of right-hand sides)
- Compute residual and obtain information on factors (optional)
- Optional restart (save data for later factorization and/or solves)
- Optional scaling (out-of-core)


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- Matrix A may be either in assembled form or a sum of element matrices
- Reverse communication interface with input by rows or by elements
- Separate calls for each phase
- Additional flexibility through user-controlled parameters (default settings minimize decisions user must make)


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- We have developed separate packages to perform these factorizations (and partial solves)
- HSL_MA54 for positive definite problems
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- Modular design helps with readability, testing, maintenance etc
- Kernels can also be reused in other solvers
- Performance can be tuned for computing environment


## Input/Output in HSL_MA77

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- analyse phase (integer data only)
- reading data for input matrix
- writing data at each node of the assembly tree
- reading data at each node
- writing reordered data ready for factorization


## Input/Output in HSL_MA77

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- writing the original real and integer data
- analyse phase (integer data only)
- factorization phase
- reading integer data at each node of the tree
- reading real data for each leaf node
- writing columns of $L$ as they are computed
- writing Schur complements to stack
- reading data from stack


## Input/Output in HSL_MA77

For HSL_MA77 to perform well, the I/O must be efficient. I/O involves:

- writing the original real and integer data
- analyse phase (integer data only)
- factorization phase
- solve phase
- reading integer/ real factor data once for forward sub. and once for back sub.


## Input/Output in Fortran

In Fortran 77/90/95 - direct access I/O is entirely record based

- Fine if every read/write is of the same amount of data
- But we need to read/write different numbers of reals and integers at each stage of the computation
- Note: we do not want to be restricted to only accessing the data in the same order as it was written so sequential access not an option


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We have got around these limitations while adhering to the strict Fortran standard by writing our own virtual memory management system

## Virtual memory management

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- The buffer is divided into fixed length pages ... a page is the same length as a record in the file
- Careful handling of the buffer within HSL_OF 01 avoids actual input-output operations whenever possible eg
- All wanted pages that are in buffer are accessed before those that are not
- When a page is freed, only written to file if it has changed


## Virtual memory management

Each set of data (such as the reals in the matrix and its factor) is accessed as a virtual array i.e. as if it were a very long array

- Long integers (64-bit) are used for addresses in the virtual array
- Most active pages of the virtual array are held in the buffer
- Any contiguous section of the virtual array may be read or written
- Each virtual array is associated with a primary file
- For very large problems, the virtual array may be too large for a single file so secondary files are used

The primary and secondary files are direct access files.

## Virtual memory management



- In this example, two superfiles associated with the in-core buffer
- First superfile has two secondaries, the second has none
- Important: user shielded from this but can control where the files are stored (primary and secondary files may be on different devices).
- Actual $\mathrm{i} / \mathrm{o}$ is not needed if user has supplied long buffer


## Use of HSL_OF 01 within HSL_MA77

- HSL_MA77 has one integer buffer and one real buffer
- The integer buffer is associated with a file that holds the integer data for the matrix $\mathbf{A}$ and the matrix factor
- The real buffer is associated with two (or three) files:
- one holds the real data for the matrix $\mathbf{A}$ and the matrix factor
- another is used for the multifrontal stack (work space)
- in the indefinite case, third file holds separate multifrontal stack for data from delayed pivots
- The user supplies pathnames together with names for the primary files


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NOTE: HSL_MA77 includes option for the files to be replaced by in-core arrays (faster if enough memory available). A combination of files and arrays may be used.

## Some numerical experiments

- Large problems from University of Florida Sparse Matrix Collection
- Double precision (64-bit) reals on single processor of a Dell Precision 670 with two 3.6 GHz Intel Xeon processors and 4 Gbytes of RAM
- f95 compiler with the -O3 option and ATLAS BLAS and LAPACK
- Comparisons with HSL solver MA57 (recall yesterday's talk)

Factor time: positive definite problems


## Solve time: positive definite problems



## Total time: positive definite problems



## Total time: indefinite problems



## Times (seconds) for large problems

| Phase | inline_1 <br> $(503,712)$ | bones10 <br> $(914,898)$ | nd24k <br> $(72,000)$ | bone010 <br> $(986,703)$ |
| :--- | :---: | :---: | :---: | :---: |
| Input | 4.87 | 6.25 | 2.86 | 8.00 |
| Ordering | 14.2 | 22.8 | 16.4 | 34.7 |
| MA77_analyse | 4.20 | 6.70 | 22.1 | 26.7 |
| MA77_factor | 90.6 | 174 | 1284 | 1491 |
| MA77_solve (1) | 5.30 | 36.0 | 10.4 | 311 |
| MA77_solve (8) | 10.6 | 41.3 | 20.7 | 331 |
| MA77_solve (64) | 60.5 | 141 | 90.2 | 499 |

MA5 7 not able to solve these on our test computer (insufficient memory).

## Unsymmetric element problems

- Also developed out-of-core multifrontal code for unsymmetric element problems. Code is called HSL_MA 78
- Based on the design of HSL_MA77
- Again uses HSL_OF 01 to handle out-of-core
- Separate kernel routine HSL_MA74 computes the partial factorization of the dense unsymmetric frontal matrices


## Comparison with frontal solver

|  | $n$ | Time (secs) |  | Factors $\left(* 10^{6}\right)$ |  |
| :--- | :---: | :---: | :---: | :---: | :---: |
|  |  | MA42_ELEMENT | MA78 | MA42_ELEMENT | MA78 |
| ship_001 | 34920 | 10.5 | 13.4 | 15.5 | 15.6 |
| m_t1 | 97578 | 552 | 101 | 135 | 56.2 |
| shipsec8 | 114919 | 950 | 101 | 196 | 55.6 |
| troll | 213453 | 3102 | 68 | 671 | 63.7 |
| fullb | 199187 | 786 | 80 | 356 | 75.1 |

These results illustrate the benefits of the multifrontal algorithm.
Appeal: We need large test problems in unassembled element form from real applications.

## Pivoting options

HSL_MA 78 offers threshold partial pivoting and threshold rook pivoting.
Threshold partial pivot: candidate must satisfy

$$
\left|f_{i j}\right| \geq \mathrm{u} * \max \left|f_{l j}\right|
$$

where $0 \leq \mathrm{u} \leq 1$ is the pivoting threshold parameter.
Threshold rook pivot: candidate must also satisfy

$$
\left|f_{i j}\right| \geq \mathrm{u} * \max \left|f_{i l}\right|
$$

ie threshold test in both row and columns.
More expensive but more stable (controls condition of $L$ and $U$ ).
Does it pay off?

## Rook versus partial pivoting

|  | $n$ | Time (secs) |  | Factors $\left(* 10^{6}\right)$ |  | Residual |  |
| :--- | ---: | ---: | :---: | :---: | :---: | :---: | :---: |
|  |  | rook | partial | rook | partial | rook | partial |
| ship_001 | 34920 | 15.0 | 13.4 | 15.6 | 15.6 | $5.7 * 10^{-16}$ | $3.1 * 10^{-16}$ |
| m_t1 | 97578 | 55.7 | 94.9 | 40.2 | 56.2 | $4.7 * 10^{-16}$ | $8.5 * 10^{-14}$ |
| shipsec5 | 179860 | $\mathbf{1 7 5}$ | 275 | 80.4 | 105 | $1.8 * 10^{-15}$ | $6.8 * 10^{-13}$ |
| ship_003 | 121178 | 146 | $\mathbf{1 1 8}$ | 70.8 | 74.0 | $7.9 * 10^{-16}$ | $1.5 * 10^{-13}$ |
| raju_001 | 151656 | 335 | $\mathbf{2 2 6}$ | 168 | 147 | $1.5 * 10^{-15}$ | $5.8 * 10^{-15}$ |

Conclude: rook pivoting can be beneficial.

## Out-of-core scaling

$$
\tilde{A}=D_{R}^{-1} A D_{C}^{-1}
$$

where $D_{R}, D_{C}$ diagonal matrices, is an equilibration of $A$ if norms of its rows and columns have approx. same magnitude.

One possibility (Ruiz)

$$
D_{R}=\operatorname{diag}\left(\sqrt{\max _{j}\left|A_{i j}\right|}\right) \text { and } D_{C}=\operatorname{diag}\left(\sqrt{\max _{i}\left|A_{i j}\right|}\right) .
$$

May be applied iteratively.
Can we implement this without explicitly assembling A?

## Out-of-core scaling

- Recall: each stage of multifrontal method involves a frontal matrix

$$
F=\left(\begin{array}{ll}
F_{1} & F_{2} \\
F_{3} & F_{4}
\end{array}\right)
$$

- $p$ rows of $F_{1}$ and $F_{2}$ are fully summed.
- $p$ columns of $F_{1}$ and $F_{3}$ are fully summed.
- Search first $p$ rows and columns of $F$ and accumulate the largest entries


## Out-of-core scaling

- Suppose row $i$ of $F$ corresponds to global row $k$ of $A$
- If $i \leq p$

$$
\left(D_{R}\right)_{k} \leftarrow \max \left(\left(D_{R}\right)_{k}, \max _{j \leq m}\left|f_{i j}\right|\right)
$$

Otherwise

$$
\left(D_{R}\right)_{k} \leftarrow \max \left(\left(D_{R}\right)_{k}, \max _{j \leq p}\left|f_{i j}\right|\right)
$$

- Similarly for $\left(D_{C}\right)_{k}$.
- Update $D_{R}$ and $D_{C}$ and then discard $F_{1}, F_{2}, F_{3}$ and stack $F_{4}$.
- Continue to next node of tree
- Avoids assembly $A$ in main memory but does require significant I/O


## Effects of equilibration

|  | Rook |  | Partial |  |
| :--- | ---: | ---: | ---: | ---: |
|  | No scaling | Scaling | No scaling | Scaling |
| x104 | 34.0 | 24.9 | 37.8 | $\mathbf{2 3 . 0}$ |
| m_t1 | 55.7 | 45.0 | 94.9 | 63.0 |
| shipsec1 | 110 | 49.3 | 174 | 44.4 |
| thread | 37.8 | 55.0 | 35.4 | 64.6 |

Notes:

- Scaling adds overhead and may not give benefit.
- But scaling can cut total cost and can be particularly beneficial for partial pivoting.
- Scaled residuals typically an order of magnitude large for partial pivoting.


## Concluding remarks

- New HSL direct solvers are performing well on large problems
- Able to solve larger problems than previously on desktop machines
- Out-of-core working adds an overhead but our memory management system attempts to minimise this (note: single rhs solve expensive)
- Scaling out of core is a new development
- Rook pivoting looks to be a useful option
- Recently we have looked at developing parallel version but out-of-core working adds complications
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Thank you and thank you John!

